

DISCRETE OPTIMIZATION WEEK 2

EXERCISE 1

We are given a finite set $S \subseteq \mathbb{R}^2$ of points in the plane, and somebody claims that these points lie approximately on a circle. How can this be verified? Here is one possible approach: An *annulus* (ring) is the region between two concentric circles. The *smallest enclosing annulus* of S is the annulus of smallest area that contains S :

If S indeed lies on a circle, then this area is 0, but also if S is “almost” on a circle, the smallest enclosing annulus will form a (small) ring around that circle.

Show that the problem of computing the smallest enclosing annulus of $S \subseteq \mathbb{R}^2$ can be formulated as a linear program!

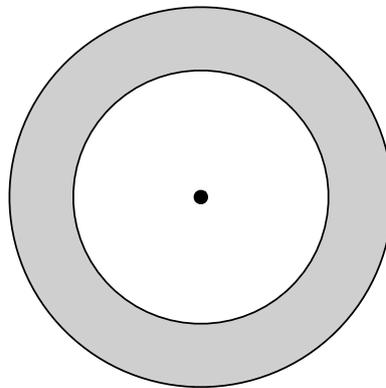


FIG. 1 – An annulus.

EXERCISE 2

Consider the following *greedy algorithm* for finding a vertex cover S in a given graph $G = (V, E)$: Set $S = \emptyset$. As long as E is nonempty, add to S a vertex v of largest degree (number of incident edges) in G . Then delete the edges incident to v from G and repeat the process with the resulting graph $G' = (V, E')$.

Show that there is no constant factor k with the property that this algorithm always finds a vertex cover whose size is at most k times as large as the size of a smallest vertex cover.

EXERCISE 3

A paper mill manufactures rolls of paper of a standard width 4 meters. But customers want to buy paper rolls of shorter width, and the mill has to cut such rolls from the 4 meters rolls.

Consider an order of

- 106 rolls of width 68 cm,
- 28 rolls of width 173 cm,
- 205 rolls of width 148 cm,
- 93 rolls of width 127 cm.

What is the smallest number of 4 m rolls that have to be cut in order to satisfy this order? Write a linear program where the solution is the answer to this question.