

## DISCRETE OPTIMIZATION WEEK 3

---

### EXERCISE 1

Consider the following *greedy algorithm* for finding a vertex cover  $S$  in a given graph  $G = (V, E)$ : Set  $S = \emptyset$ . As long as  $E$  is nonempty, add to  $S$  a vertex  $v$  of largest degree (number of incident edges) in  $G$ . Then delete the edges incident to  $v$  from  $G$  and repeat the process with the resulting graph  $G' = (V, E')$ .

Show that there is no constant factor  $k$  with the property that this algorithm always finds a vertex cover whose size is at most  $k$  times as large as the size of a smallest vertex cover.

### EXERCISE 2

We are given a set  $R$  of  $m$  red points and a set  $B$  of  $n$  blue points in the plane. We want to find a non-vertical line in the plane that passes below all red points and above all blue points (the line is also allowed to pass through points). Write a linear system of inequalities (meaning, a system of the form  $Ax \leq b$ ) that expresses whether such a line exists or not.

### EXERCISE 3

Consider the following bipartite graph :

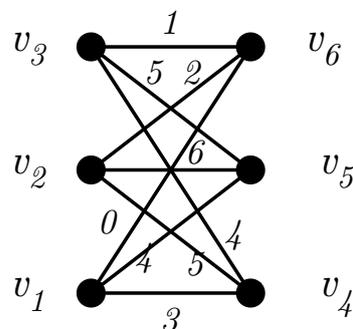


FIG. 1 –

We are interested in finding a maximum-weight perfect matching in this graph.

- Find a perfect matching of weight 10.
- Prove that there is no perfect matching of weight larger than 10, as follows : Write the linear program corresponding to the problem, as we learned in class. Multiply the six constraints corresponding to the vertices by 0, 2, 1, 3, 4, and 0, respectively ; add up these constraints ; and interpret the result.

(This is a prelude to the duality theorem that we will learn later in the course.)